

Nominal Representation of Syntax in Proof Assistants

Yee-Jian Tan

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Proof Assistant & Formal Verification

Mathematics formalisation is scaling quickly:

Fermat's Last Theorem, Navier-Stokes* | Feit-Thompson, Four colour, ...

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Life-critical software and protocols

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What makes formal verification of *programs* **challenging**?

Programming Language (PL) Meta-theory

What do we “mean” with a program?

"Semantics" meaning

$\Updownarrow$ meta-theory

"Syntax" program

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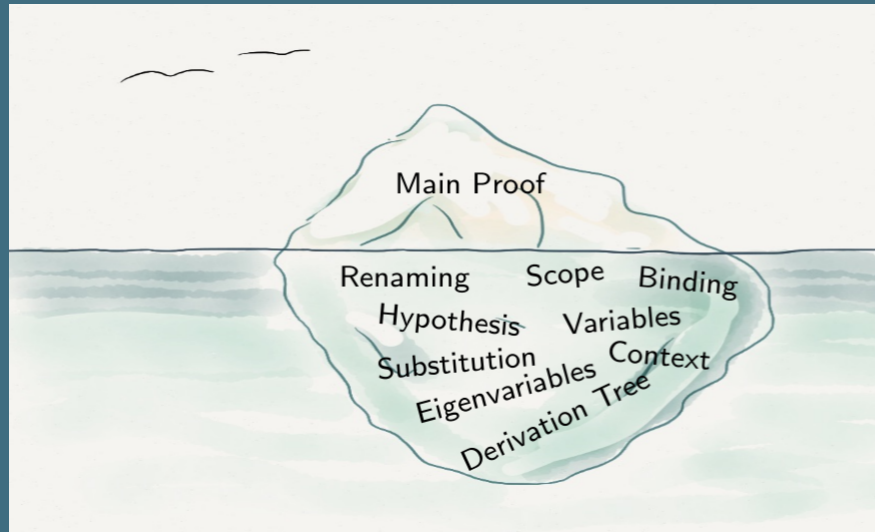
"Syntax" program

Meta-theory is **difficult**.

2005: POPLMARK: formalising PL meta-theory [Ayd+05]

2019: POPLMARK Reloaded: new benchmark problems [Abe+19]





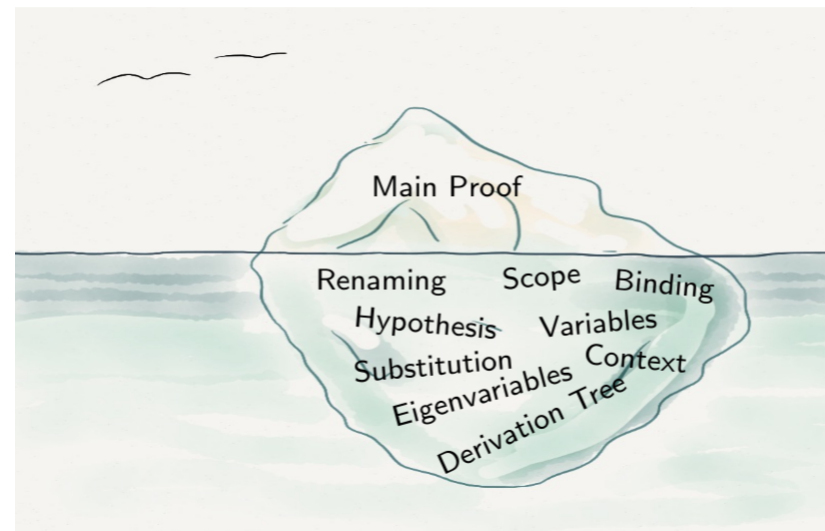
We may think of [the] proof as an iceberg. In the top of it, we find what we usually consider the real proof; underwater, the most of the matter, consisting of all mathematical preliminaries a reader must know in order to understand what is going on.

— S. Berardi [1990]

Syntax with Binders is Hard

2005: POPLMARK challenge [Ayd+05]

2019: POPLMARK Reloaded[Abe+19]



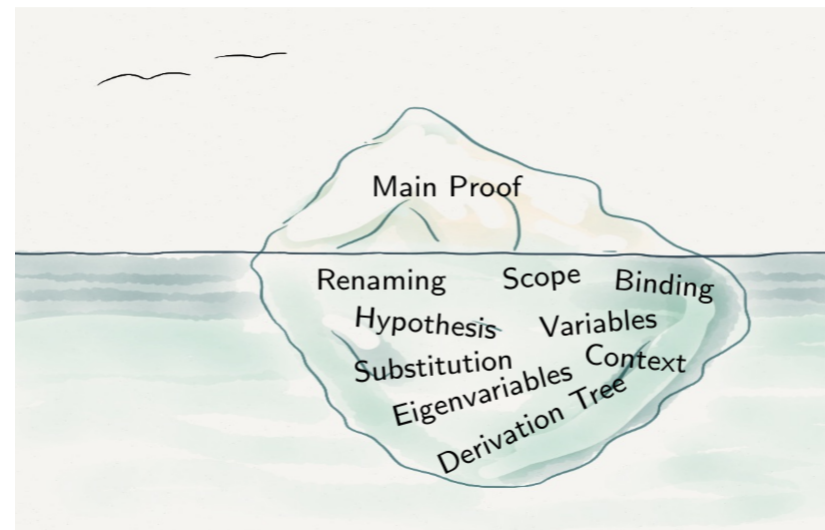
Source: “Mechanising Meta-Theory in Beluga” by
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one of the “**most critical issues**” in formalizing PL meta-theory

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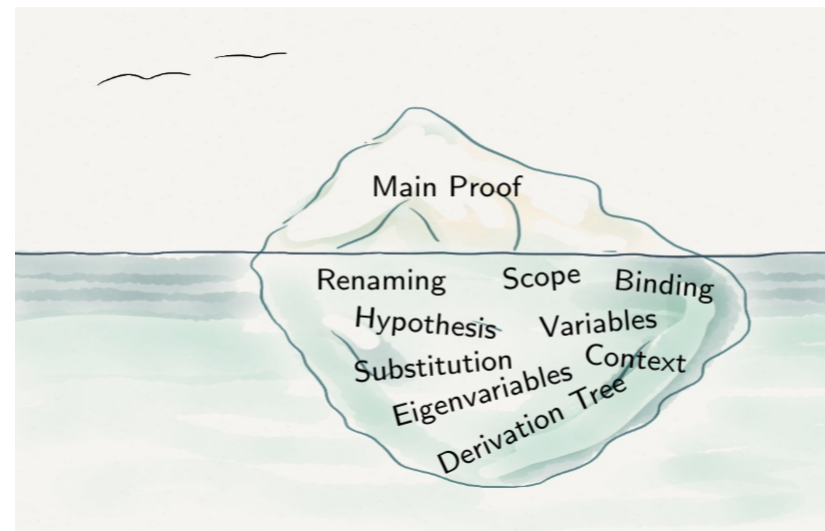
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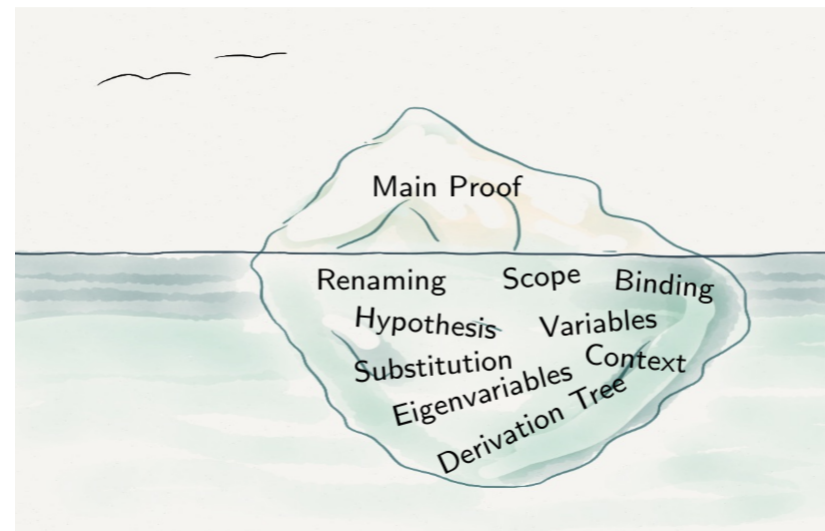
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What is a **good** variable binding representation?

Syntax with Binders: Criteria

```
1 int a;  
2 int f (int x) {  
3     return x + a;  
4 }
```

C

Syntax with Binders: Criteria

```
1 int a;  
2 int f (int x) {  
3     return x + a;  
4 }
```

C

```
1 int a;  
2 int f (int y) {  
3     return y + a;  
4 }
```

C

- α -equivalence

Syntax with Binders: Criteria

```
1 int a;  
2 int f (int x) {  
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- α -equivalence
- capture-avoidance

Syntax with Binders: Criteria

```
1 int a;  
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$\lambda.v_0 + v_1$

- α -equivalence
- capture-avoidance
- readability

Syntax with Binders: Criteria

```
1  int a;  
2  int f (int x) {  
3      return x + a;  
4  }
```

C

Induction:

Assume the binding variable is “z” ...

- α -equivalence
- capture-avoidance
- readability
- binder-induction

Syntax with binders: Comparison

	α - equivalence	capture- avoidance	readability	binder- induction	TT Impl?
Named Vars	✗	✗	✓	✗	✓
de Bruijn	✓	✓	✗	✓	✓
HOAS	✓	✓	✓	✗	✓

Syntax with binders: Comparison

	α - equivalence	capture- avoidance	readability	binder- induction	TT Impl?
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HOAS	✓	✓	✓	✗	✓
Nominal	✓	✓	✓	✓	

Syntax with binders: Comparison

	α - equivalence	capture- avoidance	readability	binder- induction	TT Impl?
Named Vars	✗	✗	✓	✗	✓
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HOAS	✓	✓	✓	✗	✓
Nominal	✓	✓	✓	✓	✗

A **Nominal TT** is needed!

Nominal TT

Breakthrough: *Nullary Parametric* TT realises **Nominal** TT [VND25].

Parametricity: “Representation Independence” [Wad89]

WP1. Implement and Evaluate **Nominal Agda**

WP2. Develop **Nominal Observational TT** (NOTT)

WP3. Implement NOTT, **unifying** with other TT variants for maximal dissemination

Yee-Jian Tan

Education

- 🇸🇬 NUS (QS 

Awards

- 🏆 *Mention de Félicitations* for L3/M1 Theses
- 🏅 Full scholarship (bachelors, masters)

Research

- 🤝 1 invited workshop talk
- 🌐 7 int'l conference/symposium (3 talks 📢)
- 📋 1 peer-reviewed publication
1 in review, 1 planned
- 🔧 4 workshop talks

Research Group

💻 DistriNet @ KU Leuven:
Hardware ↔ Software ↔ Logic

👤 D. Devriese, A. Nuyts:

- only **Prf. Assistant and (Dep)TT** subgroup in 🇧🇪
- organised **Agda Implementor's Meeting (Apr 2026)**

People

Agda Core Dev Team (A. Abel, J. Cockx)
Rocq Core Dev Team (INRIA Gallinette)
INRIA Cambium, Budapest TT Group, ...

Teaching

- 👤 Intern (Miguel Calçado, UMinho)
- 👤 Master's Thesis (Ridan Vandenberg, KU Leuven)
- 🏆 Teaching Assistant (Honour List of Student Tutors)



References

- [Ayd+05] B. E. Aydemir *et al.*, “Mechanized Metatheory for the Masses: The PoplMark Challenge,” in *Theorem Proving in Higher Order Logics, 18th International Conference, TPHOLs 2005, Oxford, UK, August 22-25, 2005, Proceedings*, J. Hurd and T. F. Melham, Eds., in Lecture Notes in Computer Science, vol. 3603. Springer, 2005, pp. 50–65. doi: 10.1007/11541868_4.
- [Abe+19] A. Abel *et al.*, “POPLMark reloaded: Mechanizing proofs by logical relations,” *Journal of Functional Programming*, vol. 29, p. e19, 2019, doi: 10.1017/S0956796819000170.
- [VND25] A. Van Muylder, A. Nuyts, and D. Devriese, “Nominal Type Theory by Nullary Internal Parametricity.” [Online]. Available: <https://arxiv.org/abs/2512.09464>
- [Wad89] P. Wadler, “Theorems for Free!,” in *FPCA '89*, Imperial College, London, United Kingdom: ACM, 1989, pp. 347–359. doi: 10.1145/99370.99404.